

10. Permanent is VNP-Complete, Part 2

Sunday, September 24, 2023 5:31 AM

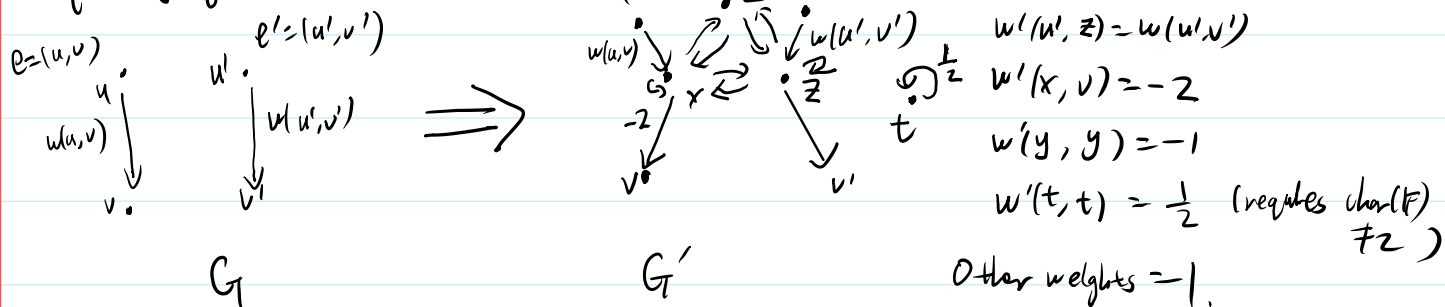
Last time: $VNP = VNP_e$.

Today: Thm (Valiant): $PERM_n = \text{perm}(X_{ij})$ is VNP-complete when $\text{char}(F) \neq 2$.

Note: If $\text{char}(F) = 2$, then $PERM = DET$ since $-1 = 1$.

(However, there are VNP-complete polynomial families even when $\text{char}(F) = 2$.)

Equality gadget:



Lemma 1: Let G be a weighted directed graph, $e, e' \in E(G)$.

Construct G' as above. Then

$$\text{perm}(G') = \sum_{C \text{ cycle cover of } G} w(C)$$

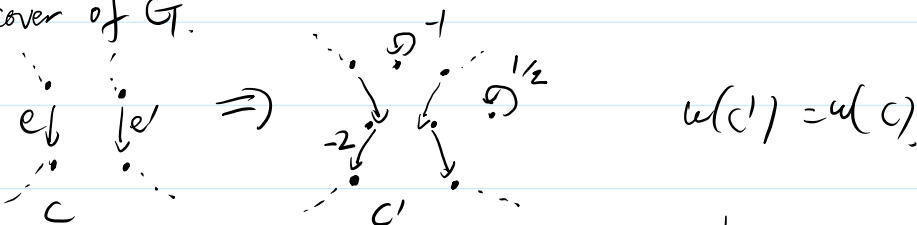
we abuse the notation and use G' to denote the weighted adjacency matrix of G' .

either $e, e' \in C$, or $e, e' \notin C$.

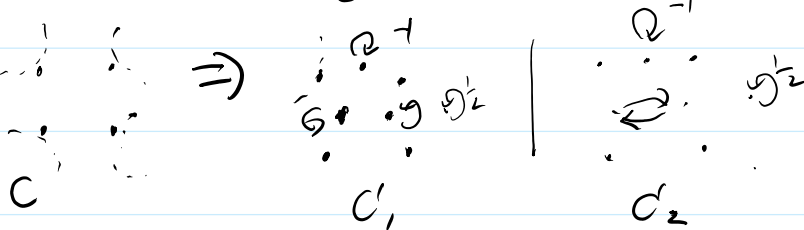
Pf: Case by case analysis.

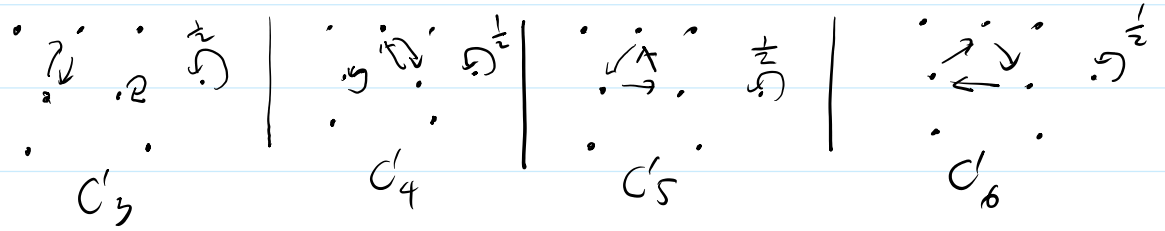
Let C be a cycle cover of G .

1) If $e, e' \in C$,



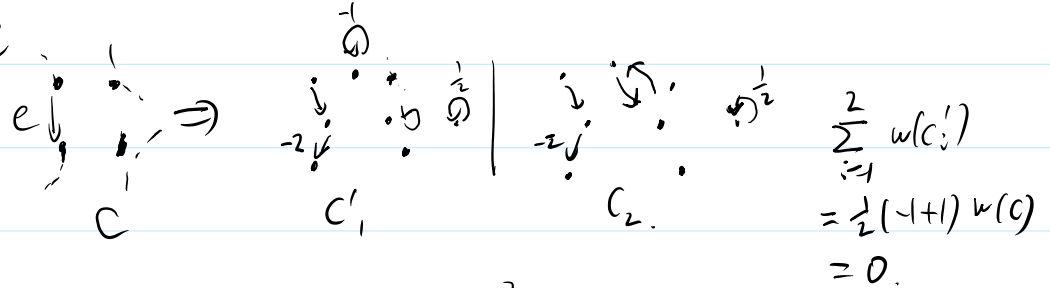
2) If $e, e' \notin C$,





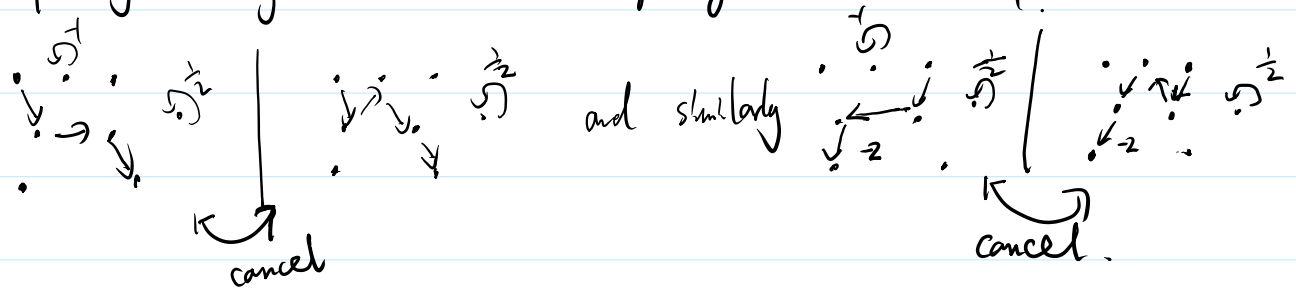
$$\sum_{i=1}^6 w(C_i) = \frac{1}{2}(-1-1+1+1+1+1) w(C) = w(C).$$

3) $e \in C, e' \notin C$



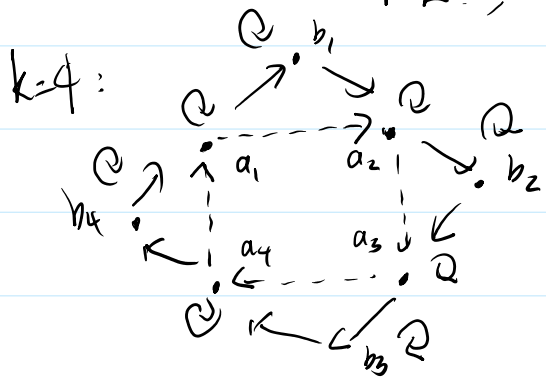
4) $e \notin C, e' \in C$. Similar to (3), $\sum_{i=1}^2 w(C_i) = 0$.

Finally, G' may have cycle covers not corresponding to those of G .



Their weights cancel each other. \square

rosette graph $R_k, k > 0$.



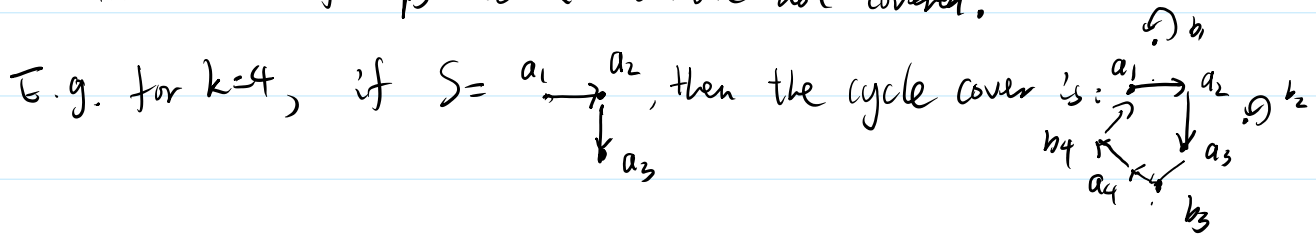
The inner edges $a_i \rightarrow a_{(i \bmod k)+1}$ are called "connector edges".

Lemma 2: Suppose S is a subset of connector edges of R_k .

1. If $S \neq \emptyset$, then \exists exactly one cycle cover of R_k that contains S and no other connector edges.

1. If $S \neq \emptyset$, then \exists exactly one cycle cover of K_k that contains S and no other connector edges.
2. If $S = \emptyset$, then \exists exactly two cycle covers of K_k that contain no connector edges.

Pf: If $S \neq \emptyset$, when $(a_i \rightarrow a_{(i \bmod k)+1}) \notin S$, use the outer edges to connect a_i and $a_{(i \bmod k)+1}$.
Then add self loops to b_i that are not covered.



It is easy to see it's the unique cycle cover.



But there is another cycle cover: consisting of only self-loops.

It is easy to see these are the only two. \square

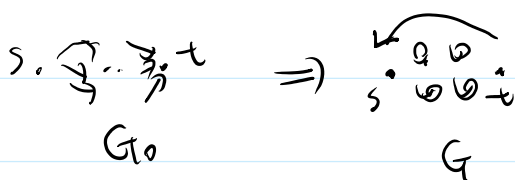
Pf of Thm: We already proved $PERM \in VMP$.

Let $(f_n) \in VMP$. Then $f_n = \sum_{e=(e_1, \dots, e_{t(n)})} g_n(x_1, \dots, x_n, e_1, \dots, e_{t(n)})$ where $(g_n) \in VP$ and t is p -bounded.

It suffices to show that $f_n \leq PERM_m$ where $m \leq \text{poly}(n)$.

As $VMP = VMP_e$ (from last time), we may assume $g_n \in VF$.

So $g_n \in VBP$, i.e., there is a $\text{poly}(n)$ -size ABP G_0 computing g_n .



The same proof that $g_n \leq p \det$ shows that $g_n = \text{perm}(G)$ (Lecture 7) where G has size $\text{poly}(n)$, and weights of G are $x_1, \dots, x_n, y_1, \dots, y_{t(n)}$.

and weights of $E(G)$ are $x_1, \dots, x_n, y_1, \dots, y_{t(n)}$, and constants in \mathbb{F} .

$$f_n = \sum_{e=(e_1, \dots, e_{t(n)}) \in \mathcal{C}(G)} g_n(x_1, \dots, x_n, e_1, \dots, e_{t(n)})$$

$$= \sum_e \sum_{\substack{\text{cycle cover} \\ C \text{ of } G}} w(C) \quad \leftarrow \text{Let } m_C := w(C), \text{ which is a monomial in } x_i\text{'s and } y_i\text{'s (with a coefficient)}$$

$$= \sum_{\substack{\text{cycle cover} \\ C \text{ of } G}} \sum_e m_C(x_1, \dots, x_n, e_1, \dots, e_{t(n)})$$

$$= \sum_{\substack{\text{cycle cover} \\ C \text{ of } G}} 2^{I(C)} m_C(x_1, \dots, x_n, 1, \dots, 1) \quad \text{where } I(C) := \# y_i \text{ not appearing in } m_C.$$

\leftarrow this is b/c if y_i appears then we must let $e_i = 1$ to get a nonzero value.

call them e_1, \dots, e_k, k_i .

If y_i does not appear, then we may let $e_i = 0$ or 1 .

Suppose k_i edges of G are labeled with y_i for $i=1, \dots, t(n)$.

Let $H_i = R_{k_i}$, where every edge has weight 1.

Let $G' =$ disjoint union of G and $H_1, \dots, H_{t(n)}$. (let $H_i = \emptyset$ if $k_i = 0$)

where the weights y_i are replaced by 1 for all i .

Finally, for each $i=1, \dots, t(n)$, and $j=1, \dots, k_i$, (call it $e'_{i,j}$) connect $e_{i,j}$ with the j -th connector edge of $H_i \cong R_{k_i}$ using the equality gadget. Call the resulting graph G'' .

$$\text{Then } \text{perm}(G'') \stackrel{\text{lem 1}}{=} \sum_{\substack{\text{cycle cover } C' \\ \text{of } G'', \text{ where} \\ \text{either } e_{i,j}, e'_{i,j} \in C' \\ \text{or } e_{i,j}, e'_{i,j} \notin C' \\ \text{for all } i,j.}} w(C') = \sum_{\substack{\text{cycle cover } C \text{ of } G \\ C_i \text{ of } H_i \\ \text{where } (C, C_1, \dots, C_{t(n)}) \\ \text{are consistent,} \\ \text{i.e. either } e_{i,j} \in C, e'_{i,j} \in C_i, \\ \text{or } e_{i,j} \notin C, e'_{i,j} \notin C_i.}} (w(C) \cdot w(C_1) \cdot \dots \cdot w(C_{t(n)}))$$

($w(C) = m_C(x_1, \dots, x_n, \vec{1})$ since we replace y_i by 1)

for all i, j .

$$= \sum_{\text{consistent}} m_C(x_1, \dots, x_n, 1, \dots, 1) \quad \left(\begin{array}{l} \text{e.g. either } e_{ij} \in C, e'_{ij} \in C, \\ \text{or } e_{ij} \notin C, e'_{ij} \notin C. \end{array} \right)$$

(b/c $w(C_i) = 1$)

$$(C_1, C_2, \dots, C_{t(n)})$$

$$= \sum_{\substack{\text{cycle cover} \\ C \text{ of } G}} m_C(x_1, \dots, x_n, 1, \dots, 1) \cdot \prod_{i=1}^{t(n)} (\# C_i \text{ consistent with } C)$$

By Lemma 2, $(\# C_i \text{ consistent with } C) = \begin{cases} 1 & Y_i \text{ appears in } m_C, \\ 2 & \text{otherwise.} \end{cases}$

So $\text{perm}(G^n) = \sum_{\substack{\text{cycle cover} \\ C \text{ of } G}} 2^{I(C)} m_C(x_1, \dots, x_n, 1, \dots, 1) = f_n. \quad \square$

Conjecture (permanent vs. determinant) (perm_n) is not a p-projection of (det_n) .

or equivalently, $\text{VBP} \neq \text{VNP}$.

(since PERM is VNP -complete, and DET is VBP -complete).

We know $\text{perm}_n \leq_p \text{det}_m$ when m is exponentially large in n .

Thm (Mignon-Ressayre '04) If $\text{perm}_n \leq_p \text{det}_m$, then $m \geq n^2/2$.